# Efficient Route Planning SS 2011

Lecture 9, Friday July 15<sup>th</sup>, 2011 (Transit Networks, GTFS)

Prof. Dr. Hannah Bast Chair of Algorithms and Data Structures Department of Computer Science University of Freiburg

### Overview of this lecture

- Organizational
  - Your feedback from Ex. Sheet #6 (transit node routing)
- Transit Networks
  - In the US, "transit" means "public transportation"
  - Transit node routing has nothing to do with this "transit"
  - We will see how to model a transit network
  - GTFS = General Transit Feed Specification
  - Do our algorithms so far work on transit networks?
  - Exercise Sheet #7: Parse a transit network from GTFS and run Dijkstra on it, and if possible, your other alg's too

# Feedback on ES#6 (transit node routing)

### Summary / excerpts Stand 15.7 12:59

- Noch beim Debuggen von contraction hierarchies
- Gerade Endsemesterstress bei vielen

BURG

INI Rell Sind jetzt auch f
ür Java ausgearbeitet

– Siehe Link auf Ihrer Daphne-Seite

https://daphne.informatik.uni-freiburg.de/svn/CodingStandards/

- Sie finden dort
  - Eine README.deutsch.txt
  - Ein vollständiges Code-Beispiel für C++
  - Ein vollständiges Code-Beispiel für Java
  - Für eigene Projekte, einfach den entsprechenden Ordner kopieren und Code schreiben, sollte dann gehen

### **Transit Networks**

What kind of data have we got?

- Stations (train stations, bus stops, etc.)
- Lines (trains, buses, trams, etc.)
- The schedule of these lines, that is, on which days do they serve which stations at which times
- Since we want to compute shortest path queries also for transit networks (best way to get from A to B), we want to model them as (directed) graphs, too, just like road networks ... but how?

Time-dependent model 1/2

- The first thing that comes to mind
  - Each station is a **node**
  - There is an arc between two nodes u and v, if there is a vehicle (train, bus, tram, ...) going non-stop from u to v
  - However, that arc can only be used at certain times, and the time it takes to travel across the arc depends on the vehicle commuting at that time
  - We can model this via a cost function for each arc (u, v)  $cost_{u,v}(t) = the time to get from u at time t ... to v$

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# Time-dependent model 2/2

#### Example

- Stations A and B with two lines L1 and L2
- L1 takes 1 hour from A to B (non-stop) and departs from A at 10:00, 14:00 and 18:00



How to compute shortest paths on such a graph?

- A simple variant of Dijkstra's algorithm does it
- Tentative distances at the nodes are now **times of day** 
  - We will store absolute times (like 10:20) and call them t[u] for node u, but we could also store times relative to the start time (like 40 minutes)
- Start with t[s] = start time and all other  $t[u] = \infty$
- When relaxing an arc (u, v) we compute  $c = c_{u,v}(t[u])$  and take t[v] = t[u] + c if that improves on the previous t[v]
- As for ordinary Dijkstra process the node u with the smallest t[u] next, and stop when this is the target node

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Time-expanded model 1/3

A node = a particular time at a particular station

- Only at times, where something (= an arrival or a departure) is happening
- For example, Freiburg Hbf @ 13:57
- There is an arc between two nodes A@t1 and B@t2 if there is a vehicle departing from A at time t1 and arriving at B at time t2, without stops inbetween
- The cost of the arc is simply the travel time  $t^2 t^1$
- There is also an arc from A@t1 to that node A@t2 with the smallest t2 > t1 ... we call these waiting arcs



How do we compute shortest paths in this model?

- It's an ordinary directed graph with non-negative arc costs, so we can use ordinary Dijkstra
- Only problem: we do not have a target node, we only have a target station
- Solution: Run Dijsktra until anyone node from the target station is settled (which will be the first one reached)

### So far, not much difference

- Given a query A@t → B, consider the sequence of arcs relaxed by a (time-dependent) Dijkstra on the timedependent graph
- The Dijkstra on the time-expanded graph relaxes the same arcs in the same order, **plus** some additional waiting arcs to some additional nodes and the arcs leaving from these nodes
- Intuitively, the time-dependent Dijkstra considers waiting and normal arcs in one (time-dependent) arc
- The big advantage of the time-expanded model is that we have an ordinary directed graph and can thus use all our previous algorithms on it

Advanced modelling issues

- For example, what about ...
  - Transfer buffers
    - We need a minimal amount of time to transfer between two vehicles → next slide
  - Service days
    - Different schedules on different weekdays, holidays, etc. → later slide
  - Multi-criteria cost functions
    - Maybe we can get from A@t to B in 3 hours with 0 transfers, or in 2 hours with 2 transfers
    - Which one is better depends on user preference, so we should compute both → next lecture

Transfer buffers 1/5

#### Time-expanded model

 This is non-trivial, because we need to distinguish between staying on a vehicle at a station (which must not require any transfer time) and changing the vehicle, for example:



Transfer buffers 2/5

Time-expanded model, solution

- Split up each node from before into an arrival node and a departure node, and add an arc between the two (we can also model layover time that way now)
- For each arrival node A@t, add a transfer node A@t' and an arc from A@t to A@t', where t' t is the transfer buffer
- For each transfer node A@t, add an arc to the departure node A@t' with the smallest t' > t  $\sqrt[n]{-700}$
- Have the waiting arcs between transfer nodes only  $D \in P$

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- Time-dependent model, solution 1
  - We also have to distinguish here between staying on a vehicle and changing the vehicle at a station
  - It looks like we can do this by simply remembering for each node, along with the tentative arrival time t[u], the id l of the vehicle with which we arrive at u
  - Then we can build the transfer buffer into the cost function

 $cost_{u,v}(t, \ell) = time to reach v, if we are at u at time t sitting in vehicle \ell$ 

 Unfortunately, Dijkstra's algorithm will not always correctly compute the shortest path anymore then



Time-dependent model, problem

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# Transfer buffers 5/5

Time-dependent model, solution 2

- Have separate arrival and departure nodes, too
- One arrival and one departure node per line suffices
- But still, we no longer only have one node per station
- Time-dependent model, solution 3
  - When we can arrive at a station at two different times t1 and t2 with different vehicles, and |t2 - t1| is ≤ the transfer buffer, pursue both possibilities
  - Then we need to do a multi-label Dijkstra (Dijkstra maintaining several shortest paths to the same node), see next lecture



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- General Transit Feed Specification
  - Standard format established by Google in 2005
  - Here is a nice story about it: <u>http://tinyurl.com/6yczek2</u>
  - See the references to the GTFS specification
  - Relatively complex, because there are so many pecularities, special cases, etc. for transit networks
  - For a simple graph model, it is easy though

### GTFS

#### Basic concepts

- stop = what we call a station
  - e.g. Freiburg Hbf or Bertoldsbrunnen
- trip = journey of a particular vehicle at a particular time
  - e.g. the journey of Bus 11 from Munzinger Straße at 9:28 to Paduaalle at 10:11
- route = trips that have a common description
  - e.g. all journeys of Bus 11 over the day
- service days = days of the week when a trip is available
  - e.g. on weekdays (Mo-Fr) or on the weekend (Sa-Su)

ZW

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#### The files we need for exercise sheet #7

- stop\_times.txt : the actual schedule information, what eventually becomes the arcs in the transit graph
- frequencies.txt : some lines repeat in exactly the same way over the same day, then you have the schedule only once in stop\_times.txt, and the periodicity here
- calendar.txt : service day patterns, and which days of the week belong to it
- trips.txt : tells us which trips commute on which service days (via the patterns from calendar.txt)
- All files are in CSV format = a table with one record per line, columns separated by a comma, headings in first line

### Write a class CsvParser

- With a method readNextLine() that reads the next line from the file and makes the columns available via another method getItem(int i)
- You find my CsvParser.h and CsvParserTest.cpp in the course SVN under folder lectures
- Note: for efficiency reasons, I let my getItem method return a const char\* which points to part of an internal string object containing the last line read

It's slightly easier if you just return an std::string, but then you get one or two additional copies / allocations

# Implementation Advice 2/7

### Graph class

If you have a graph class with members

vector<Node> \_nodes;

vector<vector<Arc> > \_adjacencyLists;

you can use that for the (time-expanded) transit network as well

- That way, you can run your algorithms with little or no modifications on the transit network as well
- You might want to add some additional info to the Node class (like the station to which a node belongs) and to the Arc class (like the name of the GTFS route)

Parse the GTFS files in this order and way

- First parse calendar.txt and remember (in a hash set) those service ids which contain the given weekday
- Then parse trips.txt and remember (in a hash set) those trip ids with a valid (= remembered) service id
- Then parse stops.txt and store (in a hash map) the names and coordinates by stop id
- Then parse frequencies.txt and store by trip id
- Then parse stop\_times.txt and for each block of lines in the file with the same trip id, add the corresponding nodes and arcs to your graph

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- Blocks with same trip id in stop\_times.txt
  - For the last step, it is very convenient to have all lines with the same trip id together in one block, and within this block have them sorted by stop\_sequence
  - You can easily achieve this with a command-line sort

sort -t, -k1,1r -k5,5n stop\_times.txt > new\_file.txt

 If you have frequency information for the trip id of a line from stop\_times.txt, repeat accordingly, for example:

### Implementation Advice 5/7

- Arrival, departure, transfer nodes ser
  - Create all nodes already while processing stop\_times.txt
  - And also the arcs from arrival to departure nodes and vice versa, and from arrival nodes to transfer nodes
  - While processing stop\_times.txt, maintain (in a hash map) for each station the list of nodes of that station, their time, and their type (arrival, departure, transfer)
  - After processing stop\_times.txt, for each station do:

Sort the nodes by time; then it is easy to add the missing arcs **from** the transfer nodes (waiting arcs to the next transfer node, and boarding arcs to the next departure node)

• Beware: several nodes with exactly the same time

DEY

AKY

Implementation Advice 6/7

Arcs from transfer nodes, example





Implementation Advice 7/7

#### Tricks to save some arcs

- We can trivially **contract** all **departure** nodes
- This replaces pairs of a boarding and a traveling arc by a single arc ... and actually **descrease** the total number of arcs in the graph, for example:

Road vs. Transit Networks

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- Assume the time-expanded model
  - Then we can all our algorithms so far also for transit networks
  - But will the speed-up over ordinary Dijsktra be the same?

### References

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### Transit network models

Timetable information: Models and Algorithms Müller-Hannemann, Schulz, Wagner, Zaroliagis, ATMOS 2007 <u>http://www.springerlink.com/content/x54715k627860283/</u>

Road Networks vs. Transit Networks

Car or Public Transport — Two Worlds

Hannah Bast, Efficient Algorithms 2009, LNCS 5760

http://www.springerlink.com/content/y46257m66372x730/

- GTFS
  - <u>http://code.google.com/transit/spec/</u> <u>transit\_feed\_specification.html</u>